

# Government College of Engineering, Jalgaon

(An Autonomous Institute of Govt. of Maharashtra)

Department of Computer Engineering

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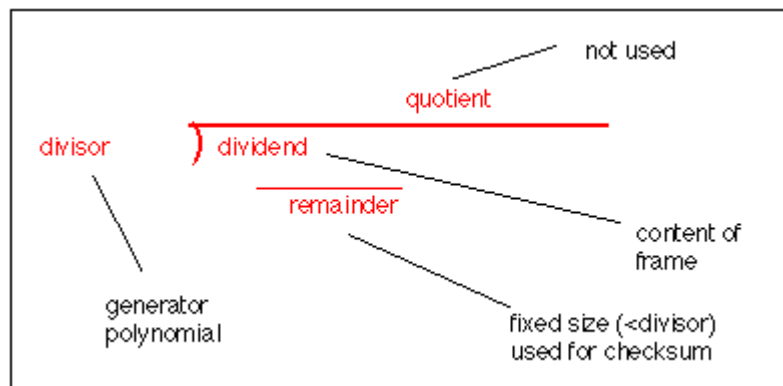
Experiment No. \_\_\_\_\_

**Objectictive:** Implementation of Cyclic Redundancy Check using modulo-2 division method.

## Theory:

A crc is a form of integrity checksum. It is a powerful method for detecting errors in the received data is by grouping the bytes of data into a block and calculating a Cyclic Redundancy Check (CRC). This is usually done by the data link protocol and calculated CRC is appended to the end of the data link layer frame.

The CRC is calculated by performing a modulo 2 division of the data by a generator polynomial and recording the remainder after division.



## Cyclic Redundancy Check

$$0 + 0 = 0$$

$$0 + 1 = 1 + 0 = 1$$

$$1 + 1 = 0 \text{ (because 2 has remainder 0 after dividing by 2)}$$

The multiplication table is equally simple:

$$0 * 0 = 0$$

$$0 * 1 = 1 * 0 = 0$$

$$1 * 1 = 1$$

What's more, subtraction is also well defined (in fact the subtraction table is identical to the addition table) and so is division (except for division by zero). What is nice, from the point of view of computer programming, is that both addition and subtraction modulo 2 are equivalent to bitwise *exclusive or* (XOR).

### Cyclic Redundancy Check and Modulo-2 Division:

- ❑ CRC or Cyclic Redundancy Check is a method of detecting accidental changes/errors in communication channel.
- ❑ CRC uses Generator Polynomial which is available on both sender and receiver side. An example generator polynomial is of the form like  $x^3 + x + 1$ . This generator polynomial represents key 1011. Another example is  $x^2 + 1$  that represents key 101.

n : Number of bits in data to be sent from sender side.

k : Number of bits in the key obtained from generator polynomial.

#### Sender Side (Generation of Encoded Data from Data and Generator Polynomial (or Key)):

1. The binary data is first augmented by adding k-1 zeros in the end of the data
2. Use *modulo-2 binary division* to divide binary data by the key and store remainder of division.
3. Append the remainder at the end of the data to form the encoded data and send the same.

**Receiver Side (Check if there are errors introduced in transmission):** Perform modulo-2 division again and if remainder is 0, then there are no errors.

### Modulo 2 Division:

The process of modulo-2 binary division is the same as the familiar division process we use for decimal numbers. Just that instead of subtraction, we use XOR here.

- In each step, a copy of the divisor (or data) is XORed with the k bits of the dividend (or key).
- The result of the XOR operation (remainder) is (n-1) bits, which is used for the next step after 1 extra bit is pulled down to make it n bits long.
- When there are no bits left to pull down, we have a result. The (n-1)-bit remainder which is appended at the sender side.

#### Example 1 (No error in transmission):

Data word to be sent - 100100

Key - 1101 [ Or generator polynomial  $x^3 + x + 1$  ]

**Sender Side:**

$$\begin{array}{r}
 \begin{array}{|c|c|}
 \hline
 1101 & 111101 \\
 \hline
 \end{array} \\
 \begin{array}{r}
 100100\ 000 \\
 \hline
 1101 \\
 \hline
 1000 \\
 1101 \\
 \hline
 1010 \\
 1101 \\
 \hline
 1110 \\
 1101 \\
 \hline
 0110 \\
 0000 \\
 \hline
 1100 \\
 1101 \\
 \hline
 001 \\
 \hline
 \end{array}
 \end{array}$$

Therefore, the remainder is 001 and hence the encoded data sent is 100100001.

**Receiver Side:**

Code word received at the receiver side 100100001

$$\begin{array}{r}
 \begin{array}{|c|c|}
 \hline
 1101 & 111101 \\
 \hline
 \end{array} \\
 \begin{array}{r}
 100100\ 001 \\
 \hline
 1101 \\
 \hline
 1000 \\
 1101 \\
 \hline
 1010 \\
 1101 \\
 \hline
 1110 \\
 1101 \\
 \hline
 0110 \\
 0000 \\
 \hline
 1101 \\
 1101 \\
 \hline
 0000 \\
 \hline
 \end{array}
 \end{array}$$

Therefore, the remainder is all zeros. Hence, the data received has no error.

**Example 2: (Error in transmission)**

Data word to be sent - 100100

Key - 1101

**Sender Side:**

$$\begin{array}{r}
 \phantom{1101} \quad 111101 \\
 \underline{1101} \quad 100100 \quad 000 \\
 \phantom{1101} \quad 1101 \\
 \phantom{1101} \quad \underline{1000} \\
 \phantom{1101} \quad \phantom{1000} \quad 1101 \\
 \phantom{1101} \quad \phantom{1000} \quad \underline{1010} \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad 1101 \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \underline{1110} \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \phantom{1110} \quad 1101 \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \phantom{1110} \quad \underline{0110} \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \phantom{1110} \quad \phantom{0110} \quad 0000 \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \phantom{1110} \quad \phantom{0110} \quad \underline{1100} \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \phantom{1110} \quad \phantom{0110} \quad \phantom{1100} \quad 1101 \\
 \phantom{1101} \quad \phantom{1000} \quad \phantom{1010} \quad \phantom{1110} \quad \phantom{0110} \quad \phantom{1100} \quad \underline{001}
 \end{array}$$

Therefore, the remainder is 001 and hence the code word sent is 100100001.

**Receiver Side:**

Let there be error in transmission media

Code word received at the receiver side - 100000001

$$\begin{array}{r}
 \phantom{1101} 111010 \\
 \hline
 1101 \overline{) 100000\ 001} \\
 \underline{1101} \phantom{000000} \\
 1010 \phantom{000000} \\
 \underline{1101} \phantom{000000} \\
 1110 \phantom{000000} \\
 \underline{1101} \phantom{000000} \\
 0110 \phantom{000000} \\
 \underline{0000} \phantom{000000} \\
 1100 \phantom{000000} \\
 \underline{1101} \phantom{000000} \\
 0011 \phantom{000000} \\
 \underline{0000} \phantom{000000} \\
 011 \phantom{000000} \\
 \hline
 \end{array}$$

Since the remainder is not all zeroes, the error is detected at the receiver side.

### Implementation:

```

//Program for CRC
#include <stdio.h>
#include <stdlib.h>
int main()
{
    int i, j, k = 0;
    int flag = 1, a[16], g[16], r[20], div[16], n, m;
    system("clear");
    printf("No. of Bite in CRC Generator : ");
    scanf("%d", &n);
    printf("\nEnter the CRC Generator : \n");
    for(i= 0; i<=n; i++)
    scanf("%d", &g[i]);
    printf("\nNo. of Bits in Data : ");
    scanf("%d", &m);
    printf("Enter Data : \n");
    for(i = 0; i<=m; i++)
    scanf("%d", &a[i]);

    if(m<n || (g[0] && g[n] == 0)
    {
        printf("Not a proper generator \n");
        exit(0);
    }
}

```

```

}

for(i = m+1; i<=m+n; i++)
    a[i] = 0;

for(j= 0; j<=n; j++)
    r[j] = a[j];

for(i= n; i<=m+n;i++)
{
    if(i>n)
    {
        for(j = 0; j<n; j++)
            r[j] = r[j+1];
        r[j] = a[i];
    }
    if(r[0])
        div[k++] = 1;
    else
    {
        div[k++] = 0;
        continue;
    }
    for(j= 0; j<=n; j++)
        r[j] = r[j] ^ g[j] ;
}

printf("\nQuotient is : ");
for(j= 0; j<k; j++)
    printf("%d ", div[j]);

printf("\nReminder is : ");
for(i= 1; i<=n; i++)
    printf("%d ", r[i]);

printf("\nTransmitted frame is : ");
for(i = m+1,j= 1; i<=m+n;i++, j++)
    a[i] = r[j];

for(i= 0; i<=m+n; i++)
    printf("%d ", a[i]);

```

```

printf("\n");

printf("\nNo. of bits of Data at Sender : ");
scanf("%d", &m);
printf("Enter Data at Sender : \n");
for(i = 0; i<=m; i++)
    scanf("%d", &a[i]);

for(j = 0; j<=n; j++)
    r[j] = a[j];
k = 0;
for(i= n; i<=m;i++)
{
    if(i>n)
    {
        for(j = 0; j<n; j++)
            r[j] = r[j+1];
        r[j] = a[i];
    }
    if(r[0])
        div[k++] = 1;
    else
    {
        div[k++] = 0;
        continue;
    }
    for(j= 0; j<=n; j++)
        r[j] = r[j] ^ g[j] ;
}

printf("\nQuotient is : ");
for(j= 0; j<k; j++)
    printf("%d ", div[j]);

printf("\nReminder is : ");
for(i= 1; i<=n; i++)
    printf("%d ", r[i]);

for(i= 1; i<=n; i++)
{
    if(r[i])
        flag = 0;
}

```

```
    }  
  
    if(flag)  
        printf("\n No Error\n");  
    else  
        printf("\nError");  
  
    return 0;  
}
```

## Output:

harish@harish-Inspiron-3537:~/harish\$ ./a.out

No. of Bite in CRC Generator : 3

Enter the CRC Generator :

1  
1  
0  
1

No. of Bits in Data : 5

Enter Data :

1  
0  
0  
1  
0  
0

Quotient is : 1 1 1 1 0 1

Reminder is : 0 0 1

Transmitted frame is : 1 0 0 1 0 0 0 0 1



No. of bits of Data at Sender : 8

Enter Data at Sender :

1  
0  
0  
1  
0  
0  
0  
0  
1

Quotient is : 1 1 1 1 0 1

Reminder is : 0 0 0

No Error

harish@harish-Inspiron-3537:~/harish\$

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